

RULES OF EQUIVALENCE for propositional logic: (Rosen same as lecture)

| | |
|-------------------------------|---|
| \rightarrow equivalence | $p \rightarrow q \equiv \neg p \vee q$ |
| \leftrightarrow equivalence | $p \leftrightarrow q \equiv (p \rightarrow q) \wedge (q \rightarrow p)$ |
| 2-negation | $\neg \neg q \equiv q$ |
| DeMorgan's Laws | $\neg(p \wedge q) \equiv \neg p \vee \neg q$ $\neg(p \vee q) \equiv \neg p \wedge \neg q$ |
| Associative ^a | $(p \wedge q) \wedge r \equiv p \wedge (q \wedge r)$ $(p \vee q) \vee r \equiv p \vee (q \vee r)$ |
| Commutative ^a | $p \wedge q \equiv q \wedge p$ $p \vee q \equiv q \vee p$ |
| Distributive | $(p \wedge q) \vee r \equiv (p \vee r) \wedge (q \vee r)$ $(p \vee q) \wedge r \equiv (p \wedge r) \vee (q \wedge r)$ |
| Idempotence | $p \vee p \equiv p$ $p \wedge p \equiv p$ |
| Inverse ^b | $p \vee \neg p \equiv \text{True}$ $p \wedge \neg p \equiv \text{False}$ |
| Identity | $p \vee \text{False} \equiv p$ $p \wedge \text{True} \equiv p$ |
| Domination | $p \wedge \text{False} \equiv \text{False}$ $p \vee \text{True} \equiv \text{True}$ |

^aOften omitted as trivial.

^bCalled "Negation" in edition 7.

Inferences rules of NATURAL DEDUCTION for propositional logic :

| | Introduction | Elimination |
|-------------------|---|--|
| AND | $\frac{p, q}{p \wedge q}$ | $\frac{p \wedge q}{p} \quad \frac{p \wedge q}{q}$ |
| OR | $\frac{p}{p \vee q} \quad \frac{q}{p \vee q}$ | $\frac{p \vee q, \begin{array}{ l} \text{from assumption} \\ p \text{ deduce } r \end{array}, \begin{array}{ l} \text{from assumption} \\ q \text{ deduce } r \end{array}}{r}$ |
| \leftrightarrow | $\frac{p \rightarrow q, q \rightarrow p}{p \leftrightarrow q}$ | $\frac{p \leftrightarrow q}{p \rightarrow q} \quad \frac{p \leftrightarrow q}{q \rightarrow p}$ |
| \rightarrow | $\frac{\boxed{\text{from assumption } p \text{ you deduce } q}}{p \rightarrow q}$ | $\frac{p, p \rightarrow q}{q}$ |
| NOT | $\frac{\boxed{\text{from assumption } p \text{ you deduce False}}}{\neg p}$ | $\frac{\neg \neg p}{p}$ |
| False | $\frac{p, \neg p}{\text{False}}$ | $\frac{\text{False}}{r}$ |

Basic additional inference rules for QUANTIFIERS:

| | | |
|-------------------|---|---|
| All-ELIM (UI) | $\frac{\forall x.P(x)}{P(c)}$ | where c is an arbitrary new constant, or one previously used |
| All-INTRO (UG) | $\frac{P(d)}{\forall x.P(x)}$ | where d is an arbitrary constant |
| Exists-ELIM (EI) | $\frac{\exists x.P(x) \quad \boxed{\text{from assumption } P(c) \text{ deduce } B}}{B}$ | where c is a new constant, which no longer occurs in B |
| Exists-INTRO (EG) | $\frac{P(b)}{\exists y.P(y)}$ | where b is any constant |

Inference rules from Rosen's textbook

Table 2: Rules of inference for Quantified Statements:

| <i>Name</i> | <i>Rule of Inference</i> |
|---------------------------------|--|
| Universal instantiation (UI) | $\frac{\forall x.P(x)}{P(c)}$ (for any arbitrary c) |
| Universal generalization (UG) | $\frac{P(c)}{\forall x.P(x)}$ for an arbitrary c |
| Existential instantiation (EI) | $\frac{\exists x.P(x)}{P(e)}$ for some element e |
| Existential generalization (EG) | $\frac{P(e)}{\exists y.P(y)}$ for some element e |

| TABLE 1 Rules of Inference. | | |
|--|--|------------------------|
| <i>Rule of Inference</i> | <i>Tautology</i> | <i>Name</i> |
| $\frac{p}{p \rightarrow q}$ $\therefore q$ | $(p \wedge (p \rightarrow q)) \rightarrow q$ | Modus ponens |
| $\frac{\neg q}{p \rightarrow q}$ $\therefore \neg p$ | $(\neg q \wedge (p \rightarrow q)) \rightarrow \neg p$ | Modus tollens |
| $\frac{p \rightarrow q}{q \rightarrow r}$ $\therefore p \rightarrow r$ | $((p \rightarrow q) \wedge (q \rightarrow r)) \rightarrow (p \rightarrow r)$ | Hypothetical syllogism |
| $\frac{p \vee q}{\neg p}$ $\therefore q$ | $((p \vee q) \wedge \neg p) \rightarrow q$ | Disjunctive syllogism |
| $\frac{p}{\therefore p \vee q}$ | $p \rightarrow (p \vee q)$ | Addition |
| $\frac{p \wedge q}{\therefore p}$ | $(p \wedge q) \rightarrow p$ | Simplification |
| $\frac{p}{q}$ $\therefore p \wedge q$ | $((p) \wedge (q)) \rightarrow (p \wedge q)$ | Conjunction |
| $\frac{p \vee q}{\neg p \vee r}$ $\therefore q \vee r$ | $((p \vee q) \wedge (\neg p \vee r)) \rightarrow (q \vee r)$ | Resolution |